# Asynchronous random sampling for decentralized detection

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# Outline

- Sequential hypothesis testing and SPRT
- Sequential change detection and CUSUM
- Decentralized detection and corresponding models
- Centralized schemes (points of reference)
- Decentralized detection using asynchronous random sampling
- Simulation comparisons

# Sequential hypothesis testing and SPRT

**Conventional binary hypothesis testing** (fixed sample size): Collection of observations $\xi_1,\ldots,\xi_K$

$$\mathbb{H}_0\colon \xi_1,\ldots,\xi_K \sim f_0(\xi_1,\ldots,\xi_K);$$
$$\mathbb{H}_1\colon \xi_1,\ldots,\xi_K \sim f_1(\xi_1,\ldots,\xi_K);$$

**Decision rule** $D(\xi_1,\ldots,\xi_K)\in\{0,1\}$

$\mathbb{P}(D=1 \mid \mathbb{H}_1)$ (Correct decision)

$\mathbb{P}(D=1 \mid \mathbb{H}_0)$ (Type I error)

$\mathbb{P}(D=0 \mid \mathbb{H}_1)$ (Type II error)

$\mathbb{P}(D=0 \mid \mathbb{H}_0)$ (Correct decision)

# Bayes and Neyman-Pearson formulation

Likelihood ratio test:

$$\frac{f_1(\xi_1,\ldots,\xi_K)}{f_0(\xi_1,\ldots,\xi_K)} \underset{\mathbb{H}_0}{\overset{\mathbb{H}_1}{\gtrless}} \gamma$$

For i.i.d.:

$$u_K = \sum_{n=1}^{K} \log\left(\frac{f_1(\xi_n)}{f_0(\xi_n)}\right) \underset{\mathbb{H}_0}{\overset{\mathbb{H}_1}{\gtrless}} \log(\gamma) = \gamma'$$

**WAIT** until $K$ samples become available, **THEN** decide

Observations $\xi_1,\ldots,\xi_n,\ldots$ are supplied **sequentially**.

$\mathbb{H}_0$: $\xi_1,\ldots,\xi_n,\ldots \sim f_0(\xi_n)$

$\mathbb{H}_1$: $\xi_1,\ldots,\xi_n,\ldots \sim f_1(\xi_n)$

| Time | Observations |
|---|---|
| 1 | $\xi_1$ |
| 2 | $\xi_1,\xi_2$ |
| ... | ... |
| $\mathcal{N}$ | $\xi_1,\ldots,\xi_{\mathcal{N}}$ |

Can $\xi_1$ make a reliable decision?

$\mathcal{N}(\xi_1,\ldots,\xi_n)=$ {stop,continue}

Time $\mathcal{N}$ is a **stopping time**

We **stop** receiving observations

**Decision Rule** $D(\xi_1,\ldots,\xi_{\mathcal{N}})\in\{0,1\}$

# WHY sequential?

For the same level of confidence with a sequential test we need, in the average, **(significantly) less samples** than a fixed sample size test, to reach a decision.

## The Sequential Probability Ratio Test (SPRT)

(Wald 1947)

Changes with time

$$u_n = \sum_{k=1}^{n} \log\left(\frac{f_1(\xi_k)}{f_0(\xi_k)}\right)$$

$$u_n = u_{n-1} + \log\left(\frac{f_1(\xi_n)}{f_0(\xi_n)}\right)$$

We define **two** thresholds $A < 0 < B$




Stopping rule: $\mathcal{N} = \inf_n \{n : u_n \notin (A, B)\}$

Decision rule: $D(\xi_1, \ldots, \xi_{\mathcal{N}}) = \begin{cases} 1 & \text{if } u_{\mathcal{N}} \geq B \\ 0 & \text{if } u_{\mathcal{N}} \leq A \end{cases}$

# Remarkable optimality property of SPRT

$$\min_{\mathcal{N},D} \mathbb{E}[\mathcal{N}|\mathbb{H}_0]$$

$$\mathbb{P}[D_{\mathcal{N}} = 1|\mathbb{H}_0] \leq \alpha; \quad \mathbb{P}[D_{\mathcal{N}} = 0|\mathbb{H}_1] \leq \beta$$

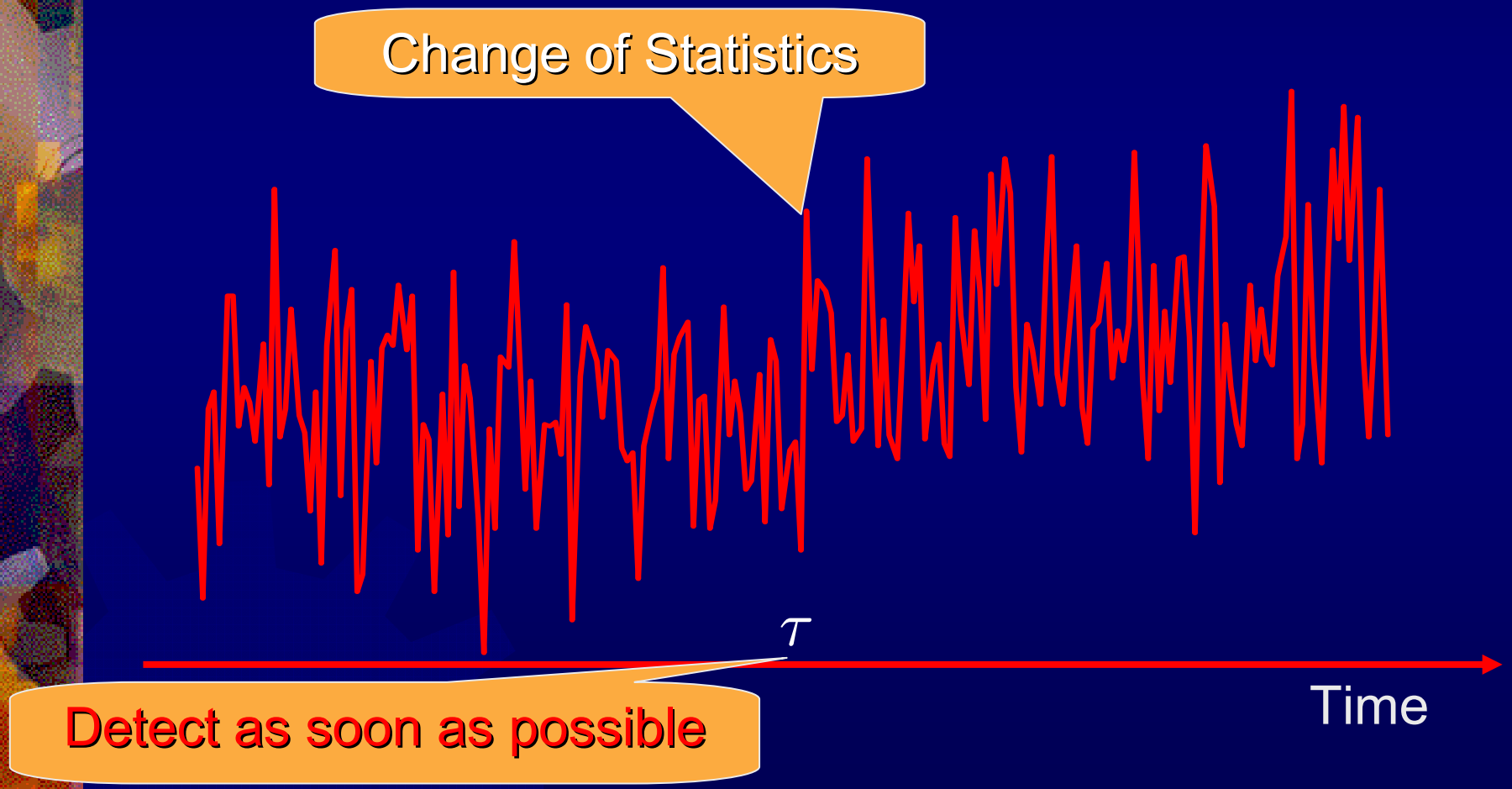
$$\min_{\mathcal{N},D} \mathbb{E}[\mathcal{N}|\mathbb{H}_1]$$

SPRT solves **BOTH** problems **simultaneously**

- Proved by Wald and Wolfowitz in 1948.

# The Sequential change detection problem

Also known as the **Disorder problem** or the **Change-Point problem** or the **Quickest Detection problem.**



# Mathematical setup

We are observing sequentially a process $\{\xi_n\}$ with the following statistics:

$$\xi_n \quad \sim f_0 \quad \text{for } 0 < n \leqslant \tau$$
$$\sim f_1 \quad \text{for } \tau < n$$

**Goal:** Detect the change time $\tau$ "**as soon as possible**"

- Change time $\tau$ : deterministic but unknown
- Densities $f_0$ , $f_1$ : known
- At every time instant $n$ we perform a test and decide whether there was a change or not. In the former case we stop in the latter we continue sampling.
- The test at time $n$ must be based on the **available information up to time $n$** (and not on any future information), i.e. **it is a stopping time.**

# Cumulative Sum (CUSUM) test

We recall the running log-likelihood:

$$u_n = \sum_{k=1}^{n} \log\left(\frac{f_1(\xi_k)}{f_0(\xi_k)}\right)$$

The running minimum: $m_n = \inf_{0 \leqslant s \leqslant n} u_s$ .

Define the CUSUM process $y_n$:

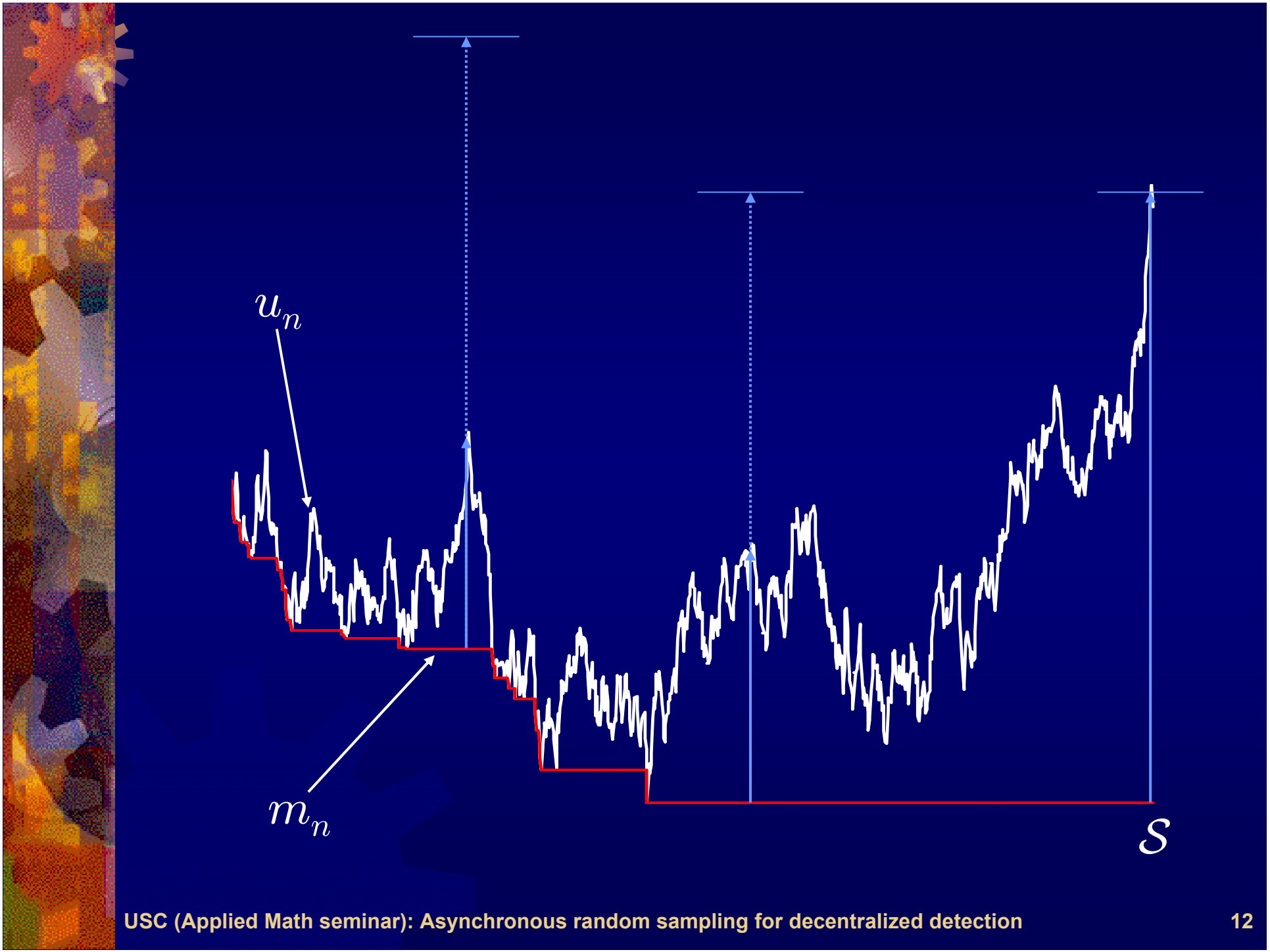
$$y_n = u_n - m_n$$

The CUSUM stopping rule:

$$\mathcal{S} = \inf_n \{\, n\colon y_n \geqslant \nu \,\}$$

We have a convenient recursion:

$$y_n = \max\left\{ y_{n-1} + \log\left(\frac{f_1(\xi_n)}{f_0(\xi_n)}\right),\ 0 \right\}$$


$u_n$
$m_n$
$\mathcal{S}$

# **Lorden's criterion** (1971)

The change time $\tau$ is deterministic and unknown.
For any stopping time $\mathcal{N}$ define the criterion:

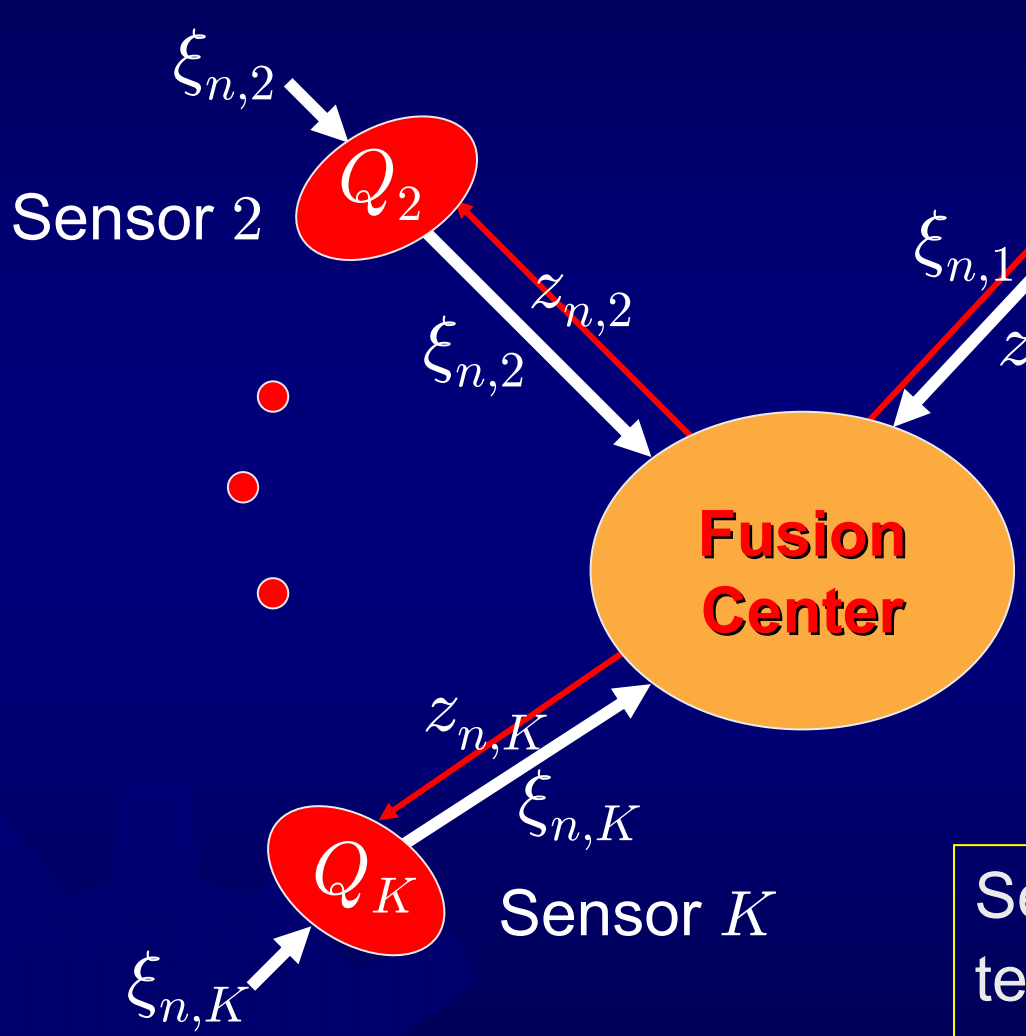
$$J_{\mathrm{L}}(\mathcal{N}) = \sup_{\tau} \operatorname{essup} \mathbb{E}_1[\ (\mathcal{N} - \tau)^+ \mid \mathcal{F}_\tau\ ]$$

**Optimization problem:** $\inf_{\mathcal{N}} J_{\mathrm{L}}(\mathcal{N})$;

subject to: $\mathbb{E}_0[\ \mathcal{N}\ ] \geqslant \gamma$ .

CUSUM solves the above optimization problem for the i.i.d. case (Moustakides 1986).

# Decentralized detection and corresponding models




**Centralized Test**
(point of reference)
High communication load

**Decentralized Test**
Quantization scheme

Sequential hypothesis testing between $f_{0,i}$ and $f_{1,i}$.

- No Local Memory:

$$z_{n,i}=Q_i(\xi_{n,i})$$

- Full Local Memory:

$$z_{n,i}=Q_i(\xi_{n,i},\xi_{n\text{-}1,i},\cdots,\xi_{1,i})$$

- Feedback with No Local Memory:

$$z_{n,i}=Q_i(\xi_{n,i},[z_{n\text{-}1,1},\cdots,z_{n\text{-}1,K}])$$

- Feedback with Partial Local Memory:

$$z_{n,i}=Q_i(\xi_{n,i},[z_{n\text{-}1,1},\cdots,z_{n\text{-}1,K}],\cdots,[z_{1,1},\cdots,z_{1,K}])$$

- Feedback with Full Local Memory:

$$z_{n,i}=Q_i(\xi_{n,i},\cdots,\xi_{1,i},[z_{n\text{-}1,1},\cdots,z_{n\text{-}1,K}],\cdots,[z_{1,1},\cdots,z_{1,K}])$$

# Centralized tests

We recall that in this case the sensors send the observations $\xi_{n,i}$ to the Fusion center.

At the Fusion center we form the running log-likelihood ratio

$$u_n = u_{n,1} + u_{n,2} + \cdots + u_{n,K}$$

$$u_{n,i} = u_{n-1,i} + \log\left(\frac{f_{1,i}(\xi_{n,i})}{f_{0,i}(\xi_{n,i})}\right)$$

and apply an SPRT:

Stopping rule: $\mathcal{N}_c = \inf_n\{n : u_n \notin (A_c, B_c)\}$

Decision rule: $D_c(u_1, \ldots, u_{\mathcal{N}_c}) = \begin{cases} 1 & \text{if } u_{\mathcal{N}_c} \geq B_c \\ 0 & \text{if } u_{\mathcal{N}_c} \leq A_c \end{cases}$

**Remark 1:**

In ALL previous detection structures it is assumed the existence of a **GLOBAL CLOCK**.
Synchronization of distant sensors with the fusion center is practically difficult (especially in sensor networks).

**Remark 2:**

In most practical applications the observation samples $\xi_{n,i}$ come from **canonical sampling of a continuous time process** $\xi_{t,i}$ where

$$\xi_{n,i} = \xi_{nT,i}$$

i.e. we sample $\xi_{t,i}$ at the time instances $t_n = nT$.

# An even better centralized scheme !

The fusion center instead of receiving the samples $\xi_{n,i}$ it can receive the **CONTINUOUS TIME PROCESSES** $\xi_{t,i}$ to form an SPRT.

$$u_t = u_{t,1} + u_{t,2} + \cdots + u_{t,K}$$

$$u_{t,i} : \text{time evolution of } u_{t,i} \text{ (sde)}$$

Stopping rule: $\mathcal{T} = \inf_t \{t : u_t \notin (A, B)\}$

Decision rule: $D(u_t, 0 \le t \le \mathcal{T}) = \begin{cases} 1 & \text{if } u_{\mathcal{T}} \ge B \\ 0 & \text{if } u_{\mathcal{T}} \le A \end{cases}$

The continuous time SPRT is **better** than the discrete time SPRT due to infinite time resolution. **It constitutes the ultimate point of reference!**

# Asynchronous random sampling

Let $t_n^i$ be increasing sequence of sampling times **NOT necessarily canonical**. At these times we sample the local log-likelihood $u_{t,i}$ in the form $u_{t_n^i,i}$ .

Instead of

$$u_t = u_{t,1} + u_{t,2} + \cdots + u_{t,K}$$

we propose the use of the following test statistic:

$$v_t = u_{t_n^1,1} + u_{t_n^2,2} + \cdots + u_{t_n^K,K}$$

Canonical sampling corresponds to: $t_n^i = nT$

Stopping rule: $\mathcal{T} = \inf_t \{t : v_t \notin (A, B)\}$

Decision rule: $D(v_1, \ldots, v_{\mathcal{T}}) = \begin{cases} 1 & \text{if } v_{\mathcal{T}} \geq B \\ 0 & \text{if } v_{\mathcal{T}} \leq A \end{cases}$

How do we transmit the local log-likelihoods $u_{t_n^i,i}$ from the sensors to the Fusion center ?

We observe

$$u_{t_n^i,i} = [u_{t_n^i,i} - u_{t_{n-1}^i,i}] + \cdots + [u_{t_1^i,i} - u_{t_0^i,i}]$$

To form the local log-likelihood $u_{t_n^i,i}$ at the fusion center, Sensor $i$ needs to transmit the differences

$$[u_{t_n^i,i} - u_{t_{n-1}^i,i}]$$

We select $t_n^i$ so that the difference $[u_{t_n^i,i} - u_{t_{n-1}^i,i}]$ takes the value $A_i$ or $B_i$ which are specified before hand.

What is the sampling strategy at Sensor $i$ ?

$$t_n^i = \inf_{t > t_{n-1}^i} \left\{ t : [u_{t,i} - u_{t_{n-1}^i,i}] \text{ hits } A_i \text{ or } B_i \right\}$$

- Every time new information arrives at the Fusion center (even from one sensor) the Fusion center updates

$$v_t = u_{t_n^1,1} + u_{t_n^2,2} + \cdots + u_{t_n^K,K}$$

and performs the test.

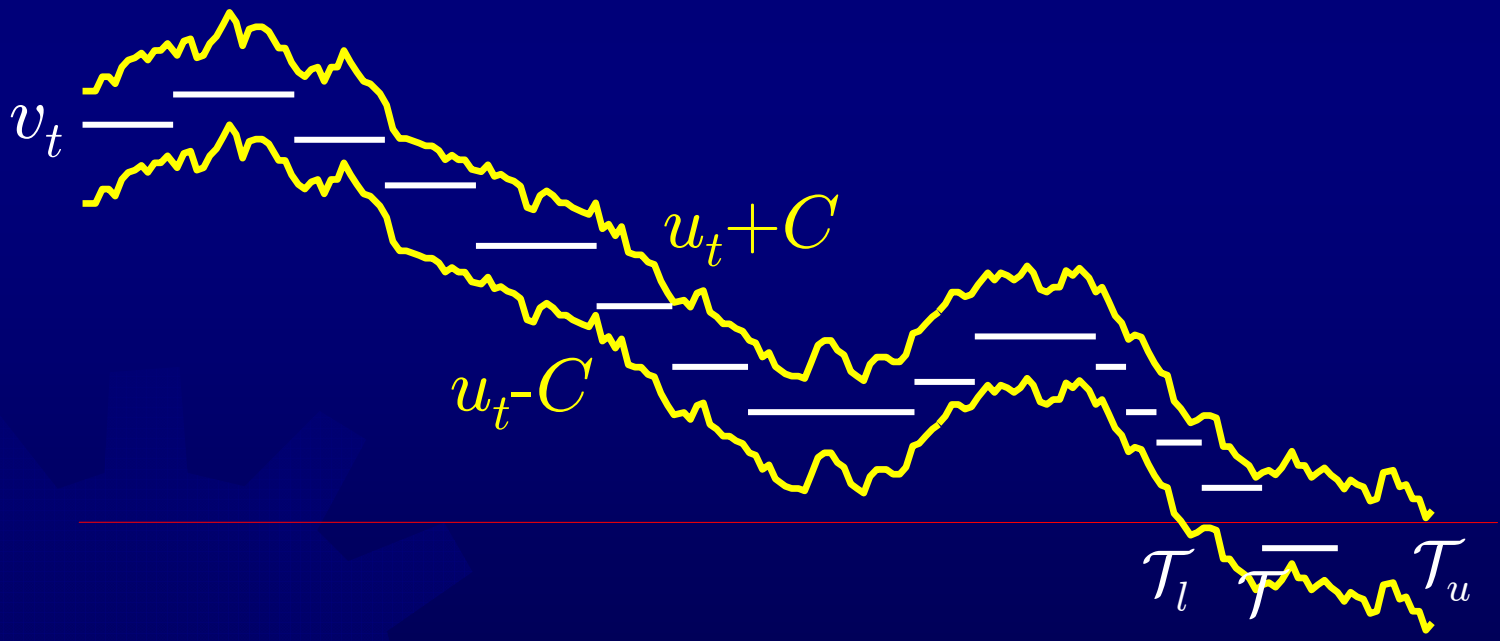
- Communication is Asynchronous and Random!!!

How do we select the local thresholds $A_i$ , $B_i$ ?

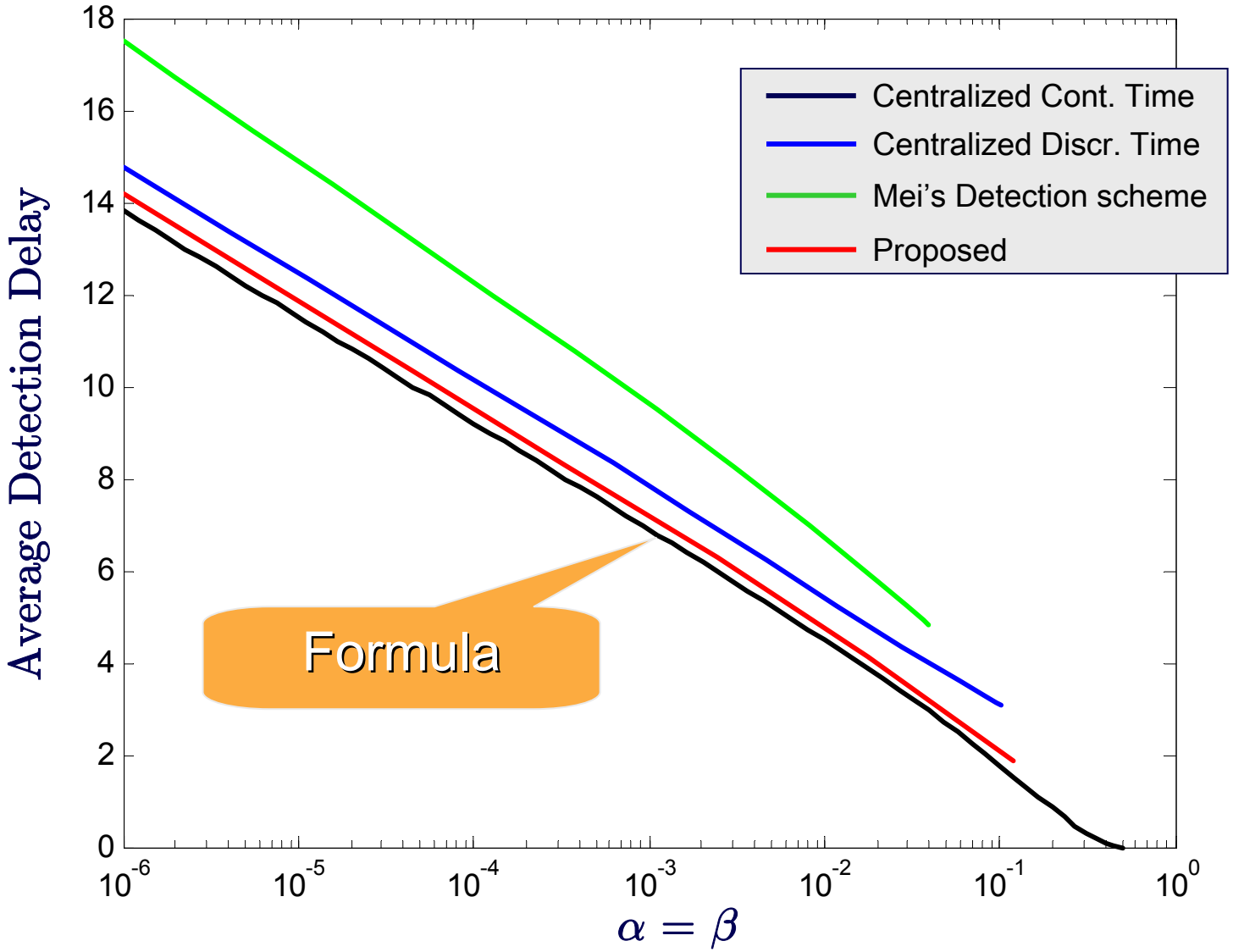
We can specify a communication rate between sensors and Fusion center.

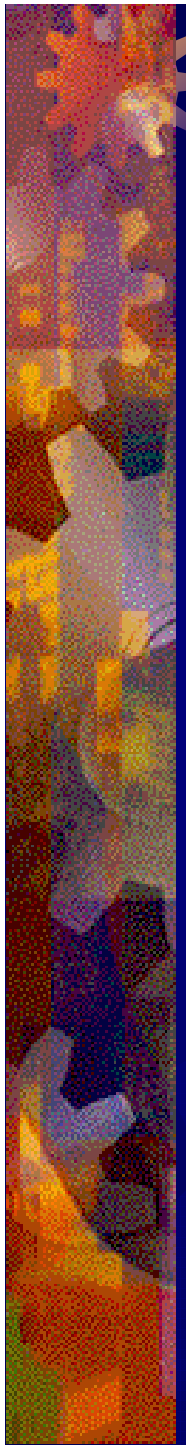
If the sensors must communicate, **in the average**, every $T$ time units, then this condition specifies completely the thresholds. We must select the thresholds so that the "average detection delay" of the local SPRTs is equal to $T$.

**Theorem:** The detection delay of the proposed scheme differs from the centralized continuous-time optimum by a constant (order-2 asymptotic optimality).



# Simulations





# END

## Thank you for your attention

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# DEPARTMENT OF MATHEMATICS

## CENTER FOR APPLIED MATHEMATICS SEMINAR


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## *"Asynchronous random sampling for decentralized detection"*


Abstract:

In decentralized detection, multiple sensors acquire information which is processed locally to produce discrete-time samples that are transmitted, at low rate, to a fusion center. The fusion center in turn collects this information from all sensors and uses it to make a sequential decision regarding a hypothesis testing or a change detection problem.

Depending on the computing capability of the sensors, the existence of feedback from the fusion center towards the sensors, we can define different models for the problem of decentralized detection. We first make a brief overview of the existing models and their optimum detection structures.

The performance of these test is then compared against the centralized test in which the fusion center has access to the complete set of information acquired by all sensors. We pay special attention to a specific class known as "sensors with full local memory" for which there was recently developed a detection scheme that approaches in performance the optimum centralized test. In the same spirit of this result we propose an alternative test that is characterized by a random sampling strategy at the sensors and a simple 1-bit asynchronous communication policy between sensors and fusion center. Our test is shown theoretically and with simulations to have performance which is very close to the performance of the optimum centralized test differing, at most, by a bounded constant.


MONDAY, 2/25/08

KAP 249

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